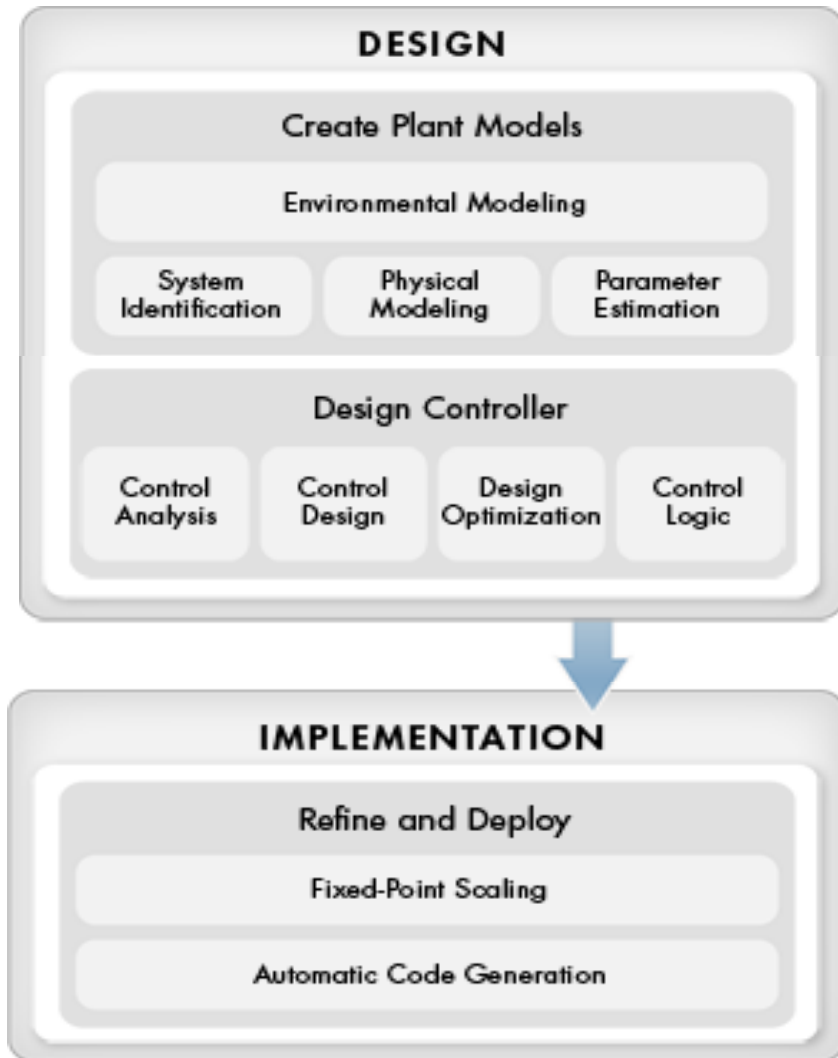


Using Heterogeneous Formal Methods in Model-Based Development

**LCCC Workshop on Formal Verification of
Embedded Control Systems**

**Bruce H. Krogh
Carnegie Mellon University in Rwanda
Kigali, Rwanda**

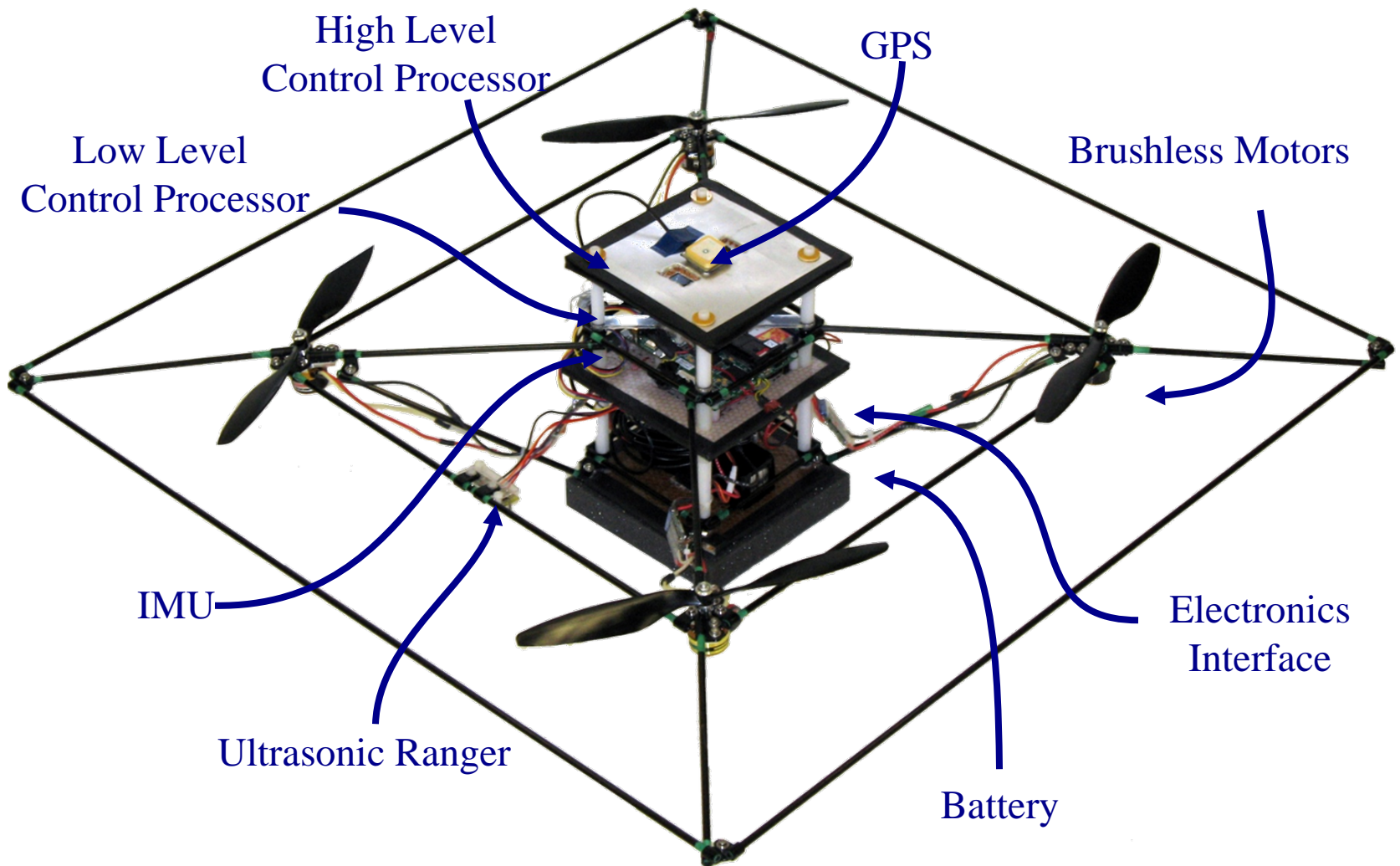
Model-Based Development



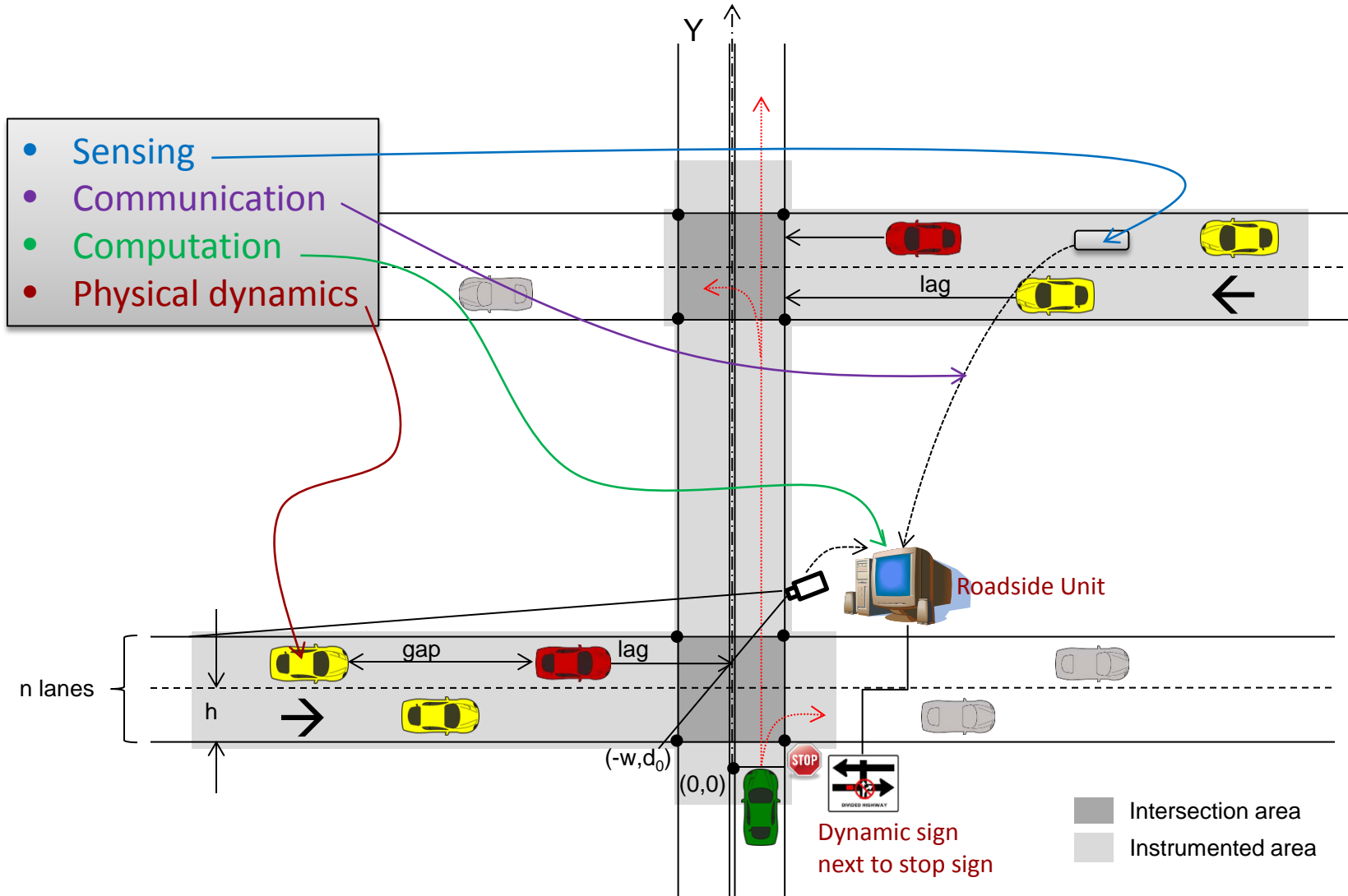
Use models to:

- **evaluate** alternative designs
- **document** the design
- **generate** implementations

Example 1: STARMAC quadrotor



Example 2: Cooperative Intersection Collision Avoidance System – Stop Sign Assist (CICAS-SSA)*



* http://www.dot.state.mn.us/guidestar/2006_2010/cicas/CICAS_SSA_ConOps_FINAL_3_18_08.pdf

Challenges in Multi-Domain MBD

No single model captures everything

- Each model represents **some** design aspect well, but not the others.
- Models are based on interdependent simplifying assumptions.
- Different **tools** focus on different properties and work only with **particular modeling formalisms**.

Challenges in Multi-Domain MBD

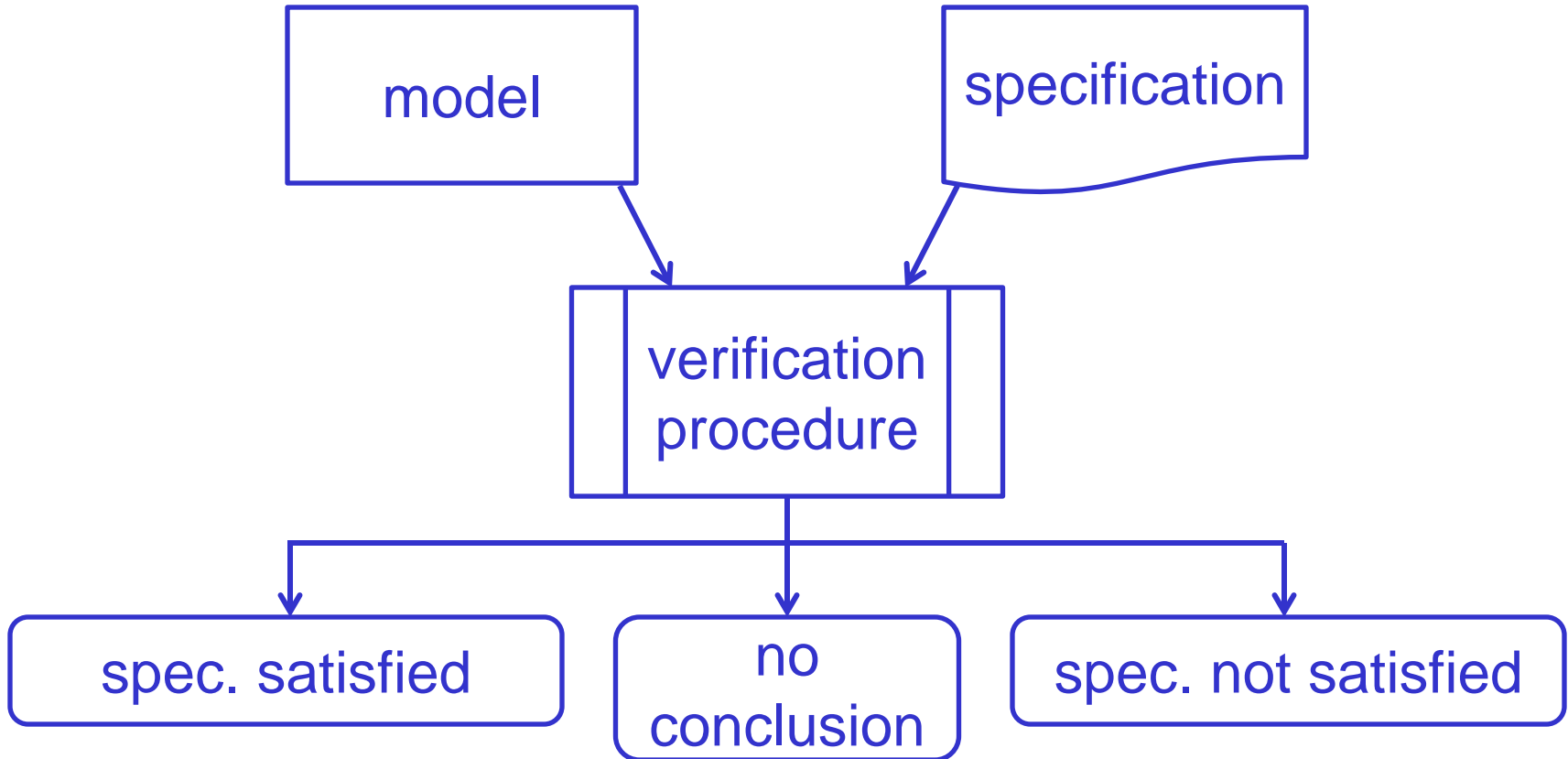
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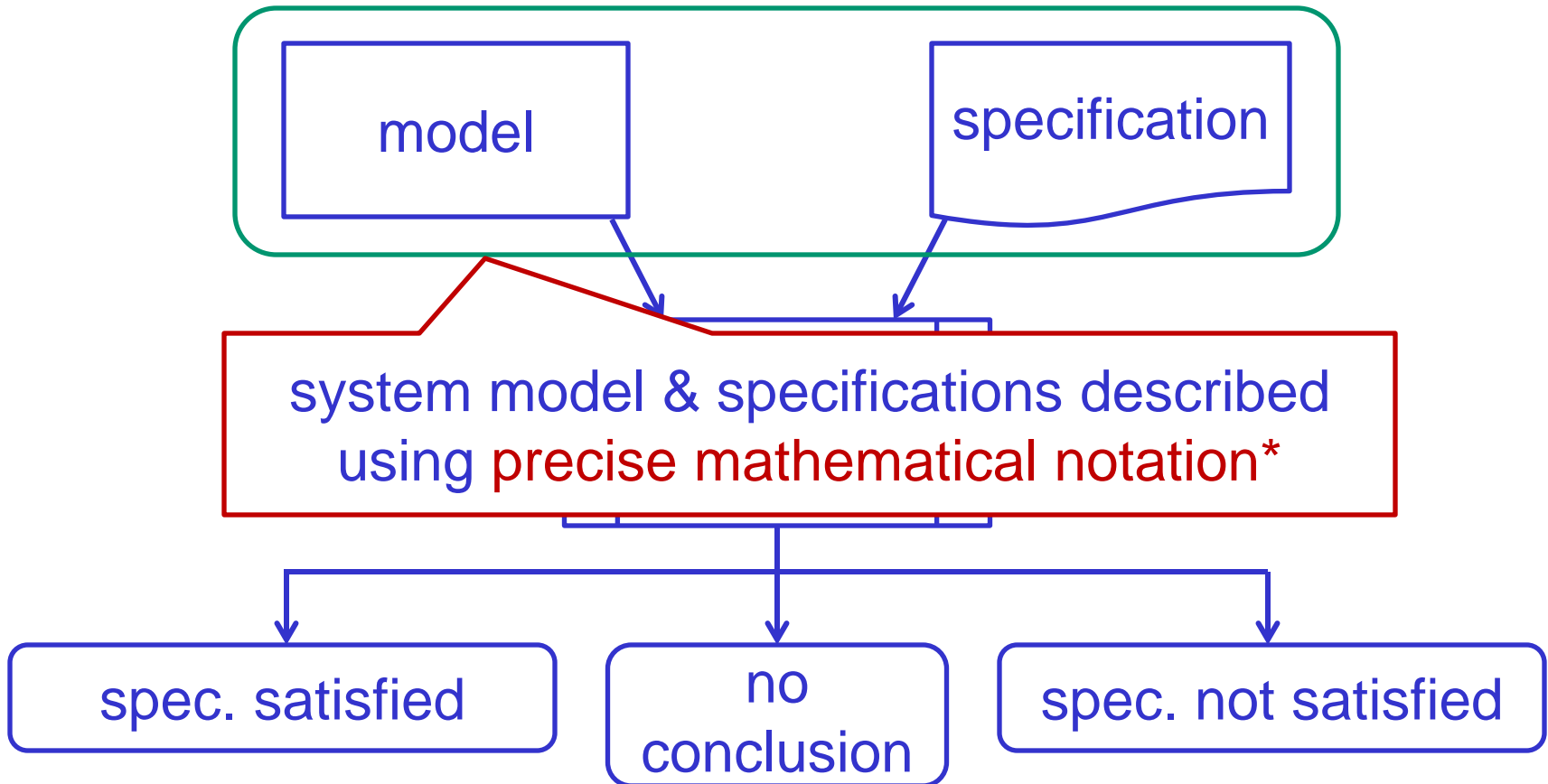
How can we:

1. Guarantee models are **consistent** with each other?
2. Infer **system-level properties** from heterogeneous analyses of heterogeneous models?

Formal Methods



Formal Methods



* unambiguous **syntax** and **semantics**

Formal Methods

verification procedure consists exclusively of well-defined mathematical operations

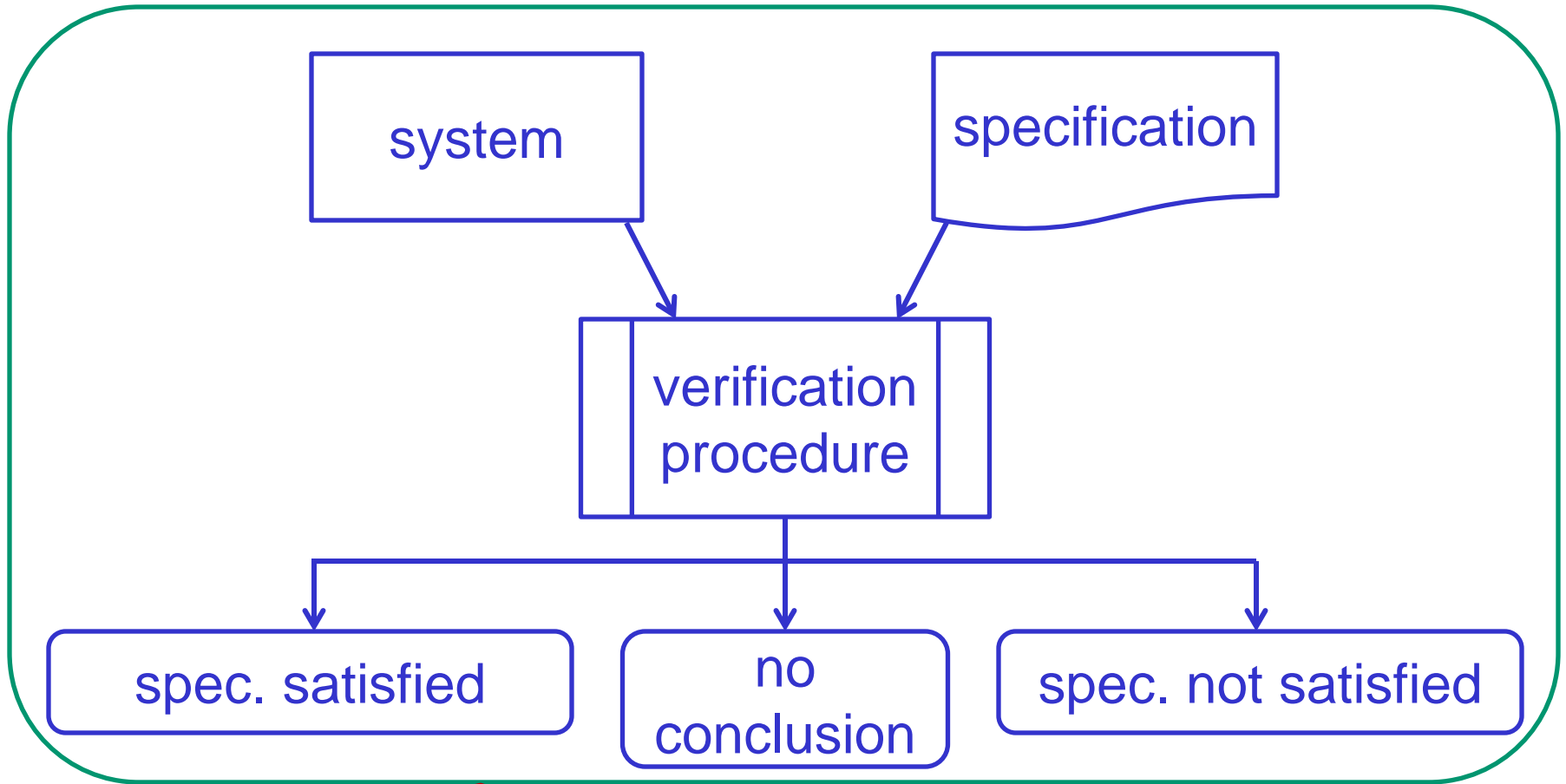
verification procedure

spec. satisfied

no
conclusion

spec. not satisfied

Formal Verification



a rigorous framework for proving specs. are satisfied

Informal Methods

system

specification

Standard Engineering Practice

- rely on experience & engineering judgment
- not provably correct
- remains the approach for most complex (and not-very-complex) systems

verification
procedure

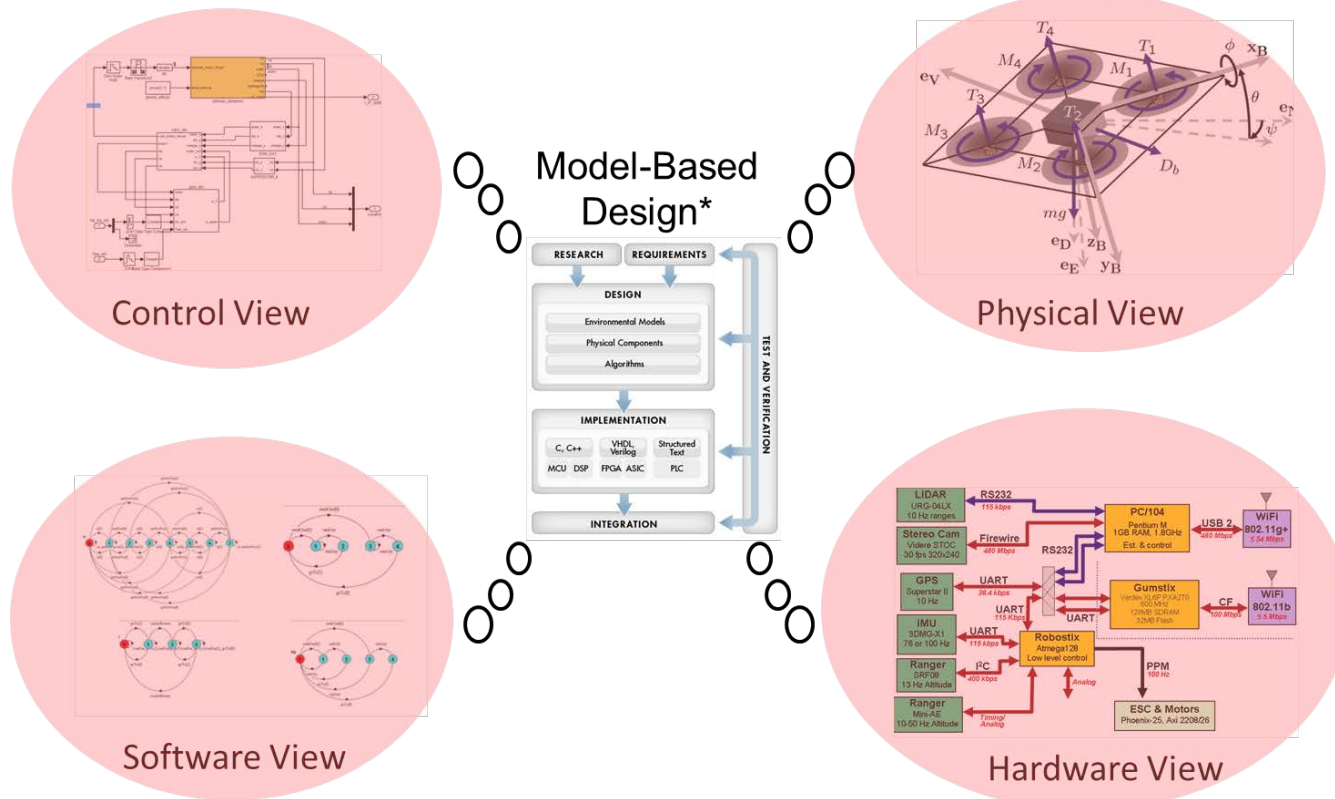
spec. satisfied

no
conclusion

spec. not satisfied

Approaches to Multi-Domain MBD:

1. Create a **universal modeling language** encompassing *everything* that needs to be modeled.



- UML/SysML (actually multiple views)
- MATLAB Simulink+Toolboxes

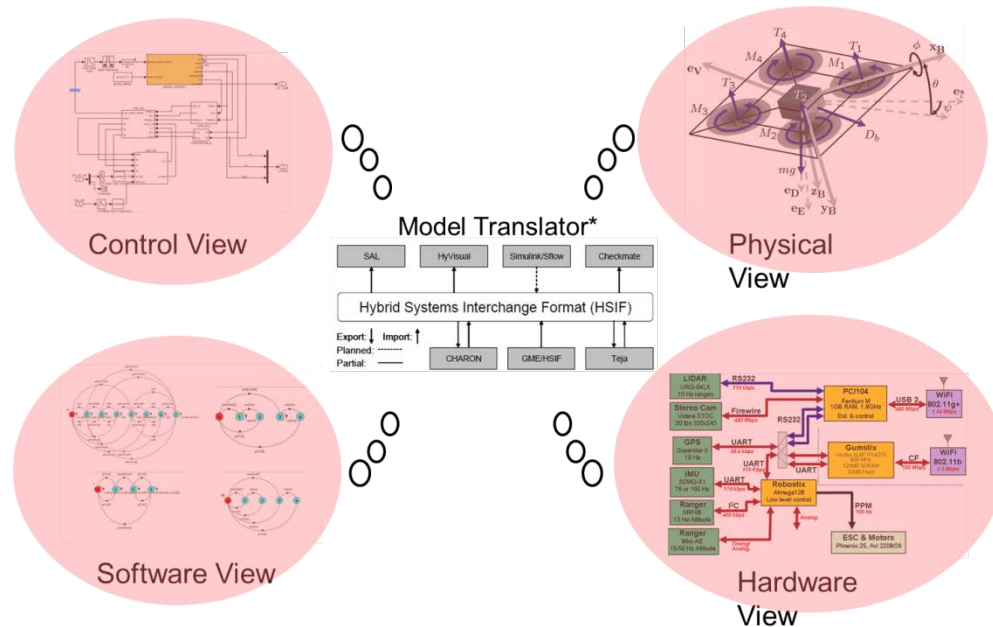
* <http://www.mathworks.com/model-based-design/>

Problems with **Universal Models**

- Comprehensive models representing *everything* become **intractable**
- Multi-domain MDB is based on **separation of concerns**: no one wants or needs the universal model
- Existing **tools operate on specific types of models**, not universal models

Approaches to Multi-Domain MBD:

2. Create tools that perform **model translation** between modeling formalisms.



- ARIES (Automatic Integration of Reusable Embedded Software)
<http://kabru.eecs.umich.edu/bin/view/Main/AIRES>
- HSIF (Hybrid Systems Interchange Format)
<http://ptolemy.eecs.berkeley.edu/projects/mobies/>

* J. Sprinkle, Generative components for hybrid systems tools, Journal of Object Technology, Mar-Apr 2003.

Problems with Model Translation

- Tool-specific translation isn't scalable
- Universal translation essentially requires a universal modeling language (Approach 1)
- Translators are difficult to maintain because modeling languages and tools continually evolve

Two Proposals for **Multi-Domain MBD**

1. How can we guarantee models are **consistent** with each other?
 - Formalize consistency at the **architectural level**
2. How can we infer **system-level** properties from heterogeneous analyses of heterogeneous models?
 - Formalize heterogeneity as **mappings between behavioral semantic domains**

(See Nikos Arechiga for point 3 in the abstract: using theorem proving to establish control design constraints)

An Architectural Framework for Multi-Domain MDB

Goal: Unify heterogeneous models through *light-weight representations* of their structure and semantics using **architecture description languages** (ADLs)

Architectural analysis:

- Does each model adhere to the base system structure & constraints (**consistency**)?
- Are all system elements represented in at least one model (**completeness**)?

Architectures for CPS

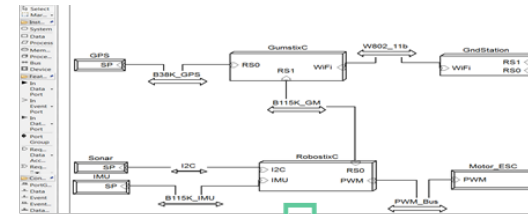
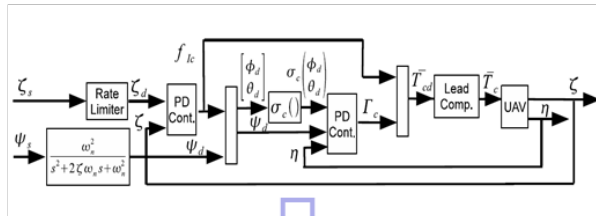
Architecture: The **set of structures** needed to **reason about the system**, which comprise functional **elements**, **relations** among them and **properties** of both.*

- CPS *base architecture* defines
 - component connectivity & physical coupling
 - data, control, & physical signal flows
- Model architectures define
 - components and connectors exposing the model structure for evaluation vis-à-vis the base architecture

* *Documenting Software Architecture: Views and Beyond*, 2nd Ed. Clements et al. 2010.

Models as Architectural Views

Model M_x



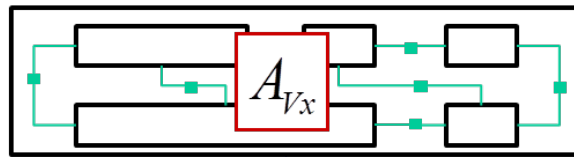
Model M_y

$$R_{V_x}^{M_x}$$

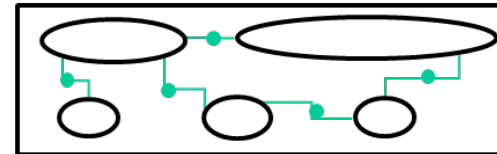
model-to-view relations

$$R_{V_y}^{M_y}$$

View V_x



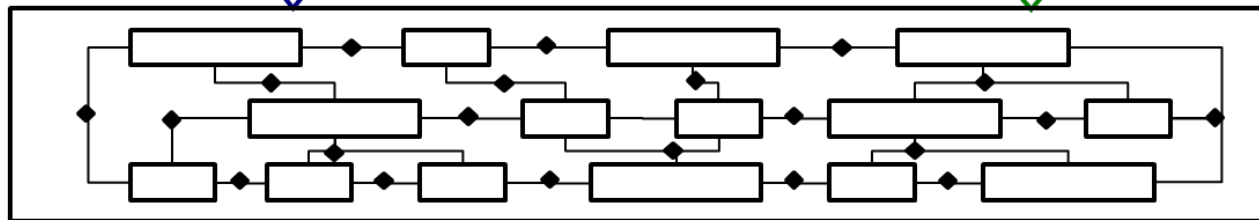
View V_y



$$R_{BA}^{V_x}$$

view-to-base-arch. relations

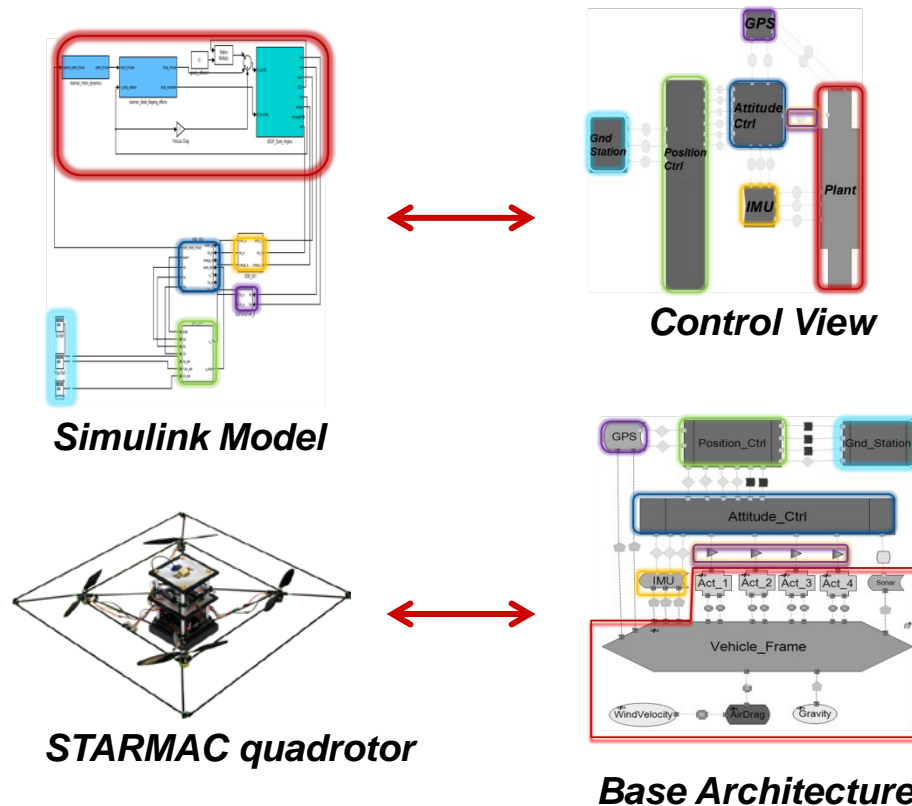
$$R_{BA}^{V_y}$$



Base CPS Architecture

Example: Consistency Analysis

- Typed-graph morphisms expose inconsistencies between model architectures and the base architecture.



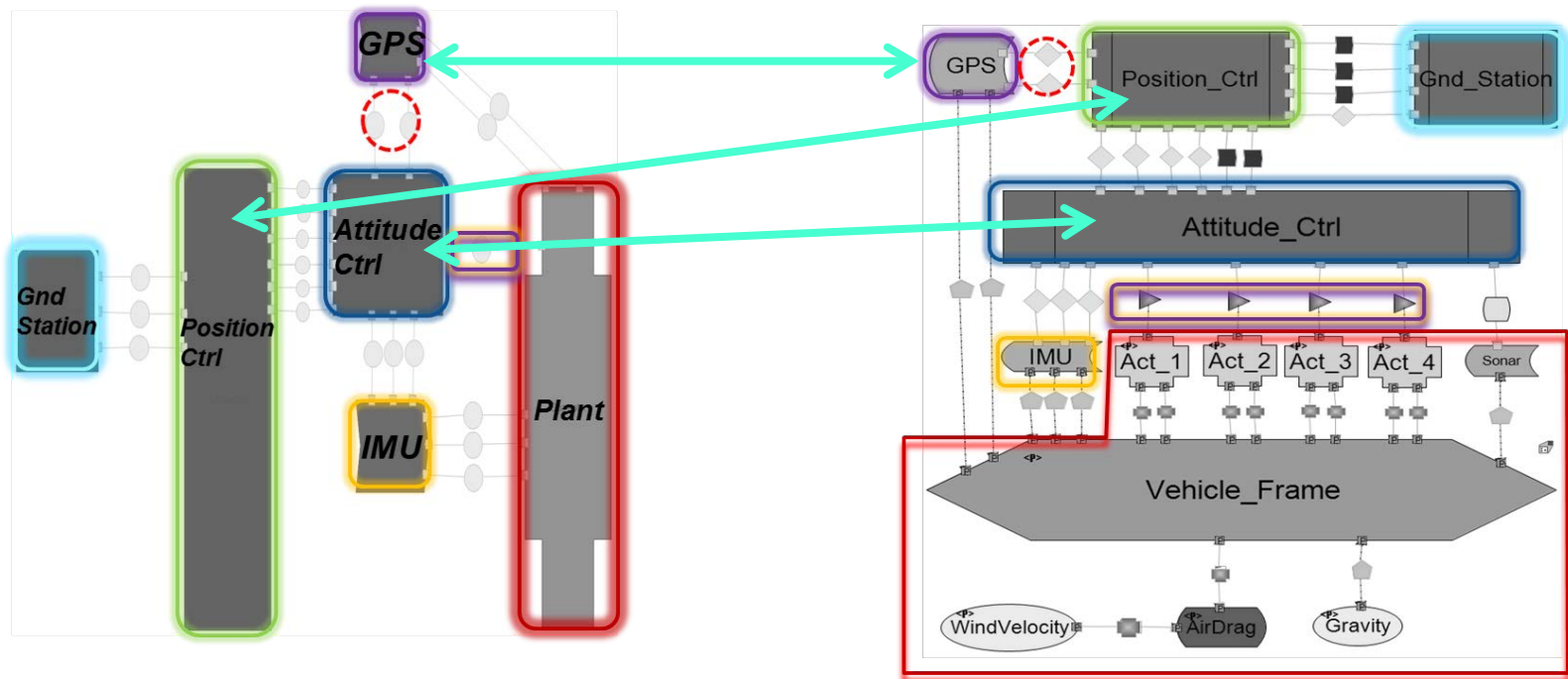
Component and connector relations defined by architecture types and semantics.

Example: Consistency Analysis

GPS sensor connected to:

attitude controller

position controller



Control View

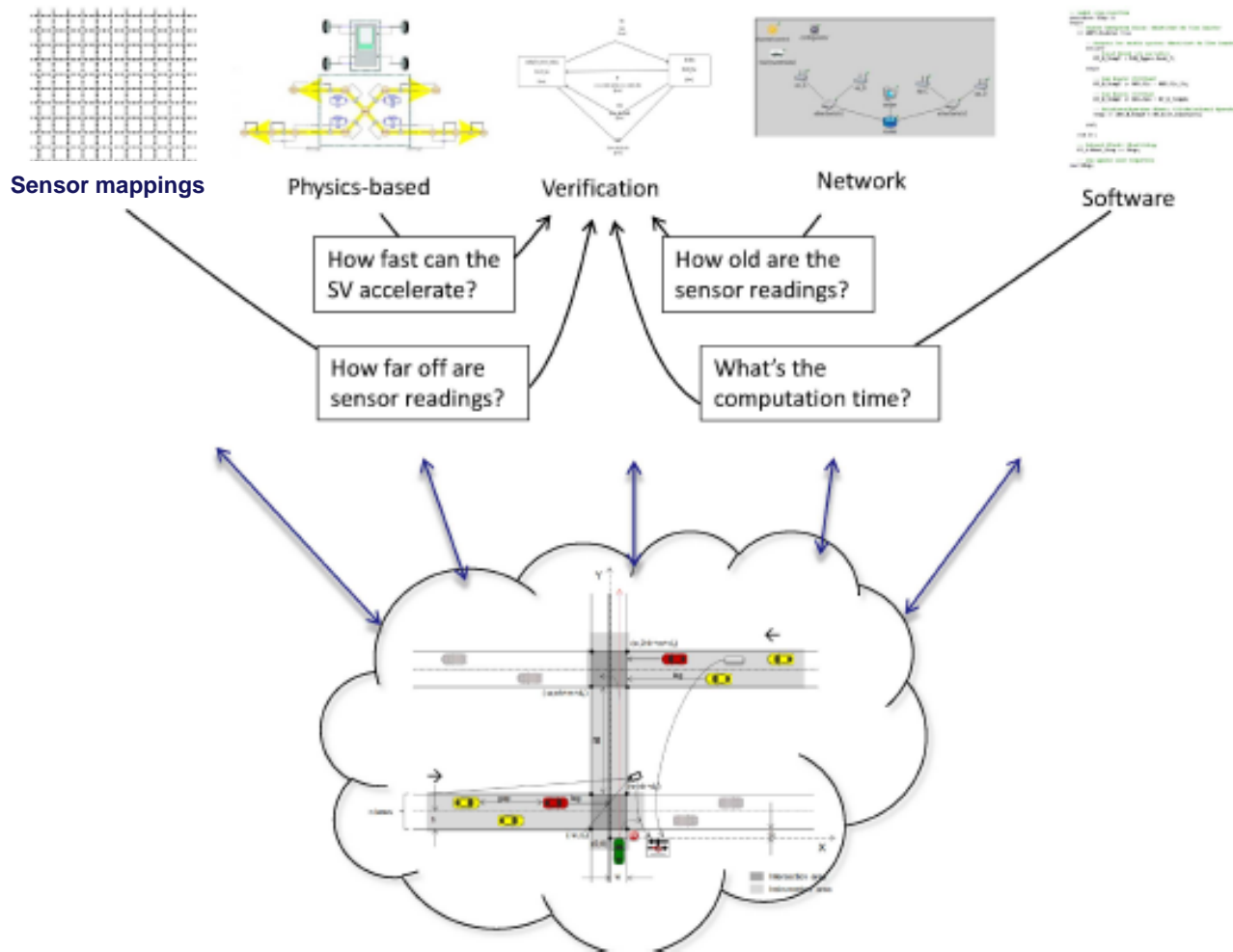
Base Architecture

Architectures for Multi-Domain MBP

- **CPS architectures**: extension of software/hardware to include physical domains
- **consistency and completeness analysis**
- **parametric consistency analysis**
- **implementation in ADL tool (ACME)**

papers: **google** “Bhave Garlan Krogh Architecture”

Heterogeneous Verification



Using **Behavioral Semantics** for Heterogeneous Verification

Goal: Provide a formal framework to perform verification using a heterogeneous set of modeling formalisms and analysis tools.

Heterogeneous Verification:

- **abstraction** using different formalisms
- **implication** using different formalisms
- **composition** using different formalisms

Using Behavioral Semantics for Heterogeneous Verification

Goal: Provide a formal framework to perform verification using a heterogeneous set of modeling formalisms and analysis techniques

Just presented last week
at HSCC 2013 by
Akshay Rajhans

- abstraction using different formalisms
- implication using different formalisms
- composition using different formalisms

Heterogeneous Verification via Behavioral Semantics

Goal: Provide a formal framework to

analyze heterogeneous

Today's presentation.

tools

is

Heterogeneous Verification:

- **abstraction** using different formalisms
- **implication** using different formalisms
- **composition** using different formalisms

Models, Specifications and Behaviors

Models: $M \in \mathcal{M}$ means the **model** M is constructed using a **modeling formalism** \mathcal{M} (e.g., state equations, Petri nets, block diagrams).

Specifications: $S \in \mathcal{S}$ means the **specification** S is constructed using a **specification formalism** \mathcal{S} (e.g., inequalities, logical expressions, automata, differential inclusions).

Behaviors: $B \subseteq \mathcal{B}$ means the **behavior domain** B is in the **class of behaviors** \mathcal{B} (e.g., traces, piecewise continuous functions, real numbers).

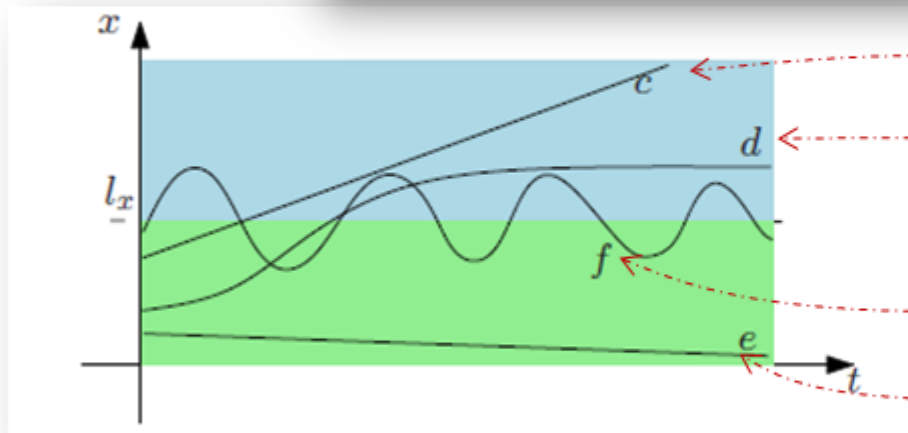
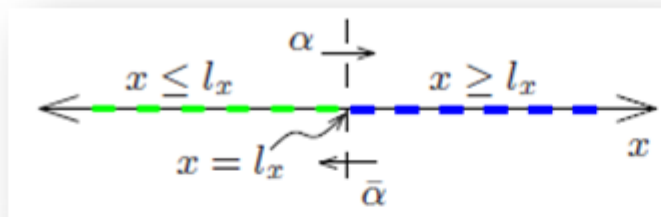
Homogeneous Behavioral Semantics

- Given a model M and a behavior domain B , the behavioral semantics for M is a set of behaviors in B , denoted $[[M]]^B$.
- Given a model S and a behavior domain B , the behavioral semantics for S is a set of behaviors in B , denoted $[[S]]^B$.
 - M_1 an **abstraction** M_0 : $[[M_0]]^B \subseteq [[M_1]]^B$.
 - S_1 **implies** S_0 : $[[S_1]]^B \subseteq [[S_0]]^B$
 - S is true for M (**entailment**) : $[[M]]^B \subseteq [[S]]^B$

Heterogeneity: Relations on Pairs of Behavioral Domains

$$R_1 \subseteq B_0 \times B_1$$

Example



\mathcal{B}_0 : continuous trajectories

B_0 : 1-d continuous trajectories in x

α

$\alpha\bar{\alpha}\alpha\bar{\alpha}\alpha\bar{\alpha}\dots$

ε

\mathcal{B}_1 : event traces

$B_1 = \{\alpha, \bar{\alpha}\}^* \cup \{\alpha b, \bar{\alpha}\}^\omega$

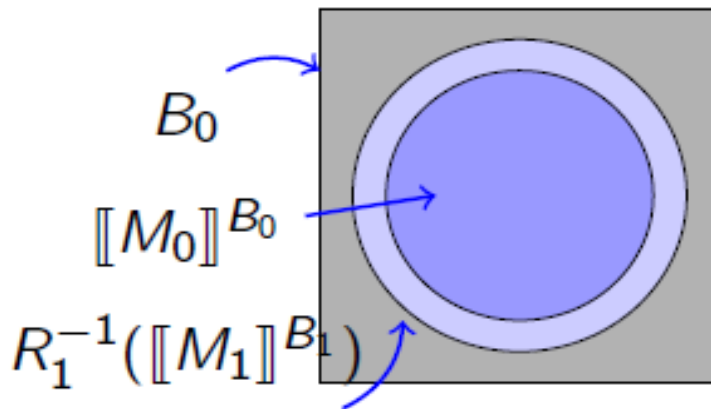
Heterogeneous Abstraction

$$R_1 \subseteq B_0 \times B_1$$

Heterogeneous Abstraction

$M_0 \sqsubseteq^{R_1} M_1$, if

$$\llbracket M_0 \rrbracket^{B_0} \subseteq R_1^{-1}(\llbracket M_1 \rrbracket^{B_1}).$$



M_1 abstracts M_0
(via R_1)

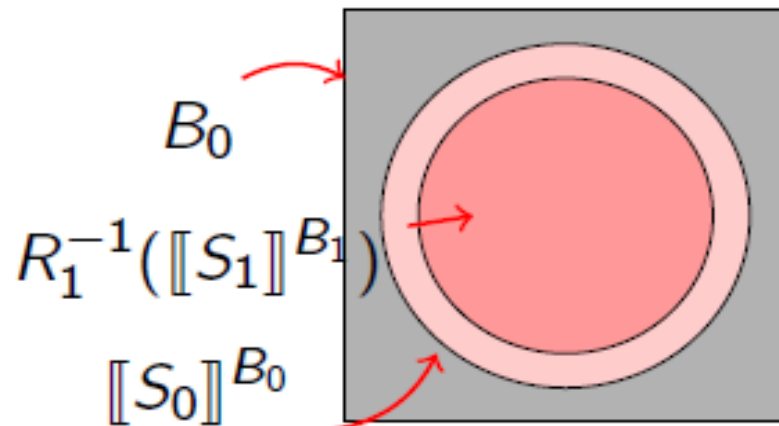
Heterogeneous Specification Implication

$$R_1 \subseteq B_0 \times B_1$$

Heterogeneous Specification Implication

$S_1 \Rightarrow^{R_1} S_0$, if

$$R_1^{-1}(\llbracket S_1 \rrbracket^{B_1}) \subseteq \llbracket S_0 \rrbracket^{B_0}.$$



S_1 stronger than S_0
(via R_1)


Heterogeneous Verification

$$R_1 \subseteq B_0 \times B_1$$

Heterogeneous Verification

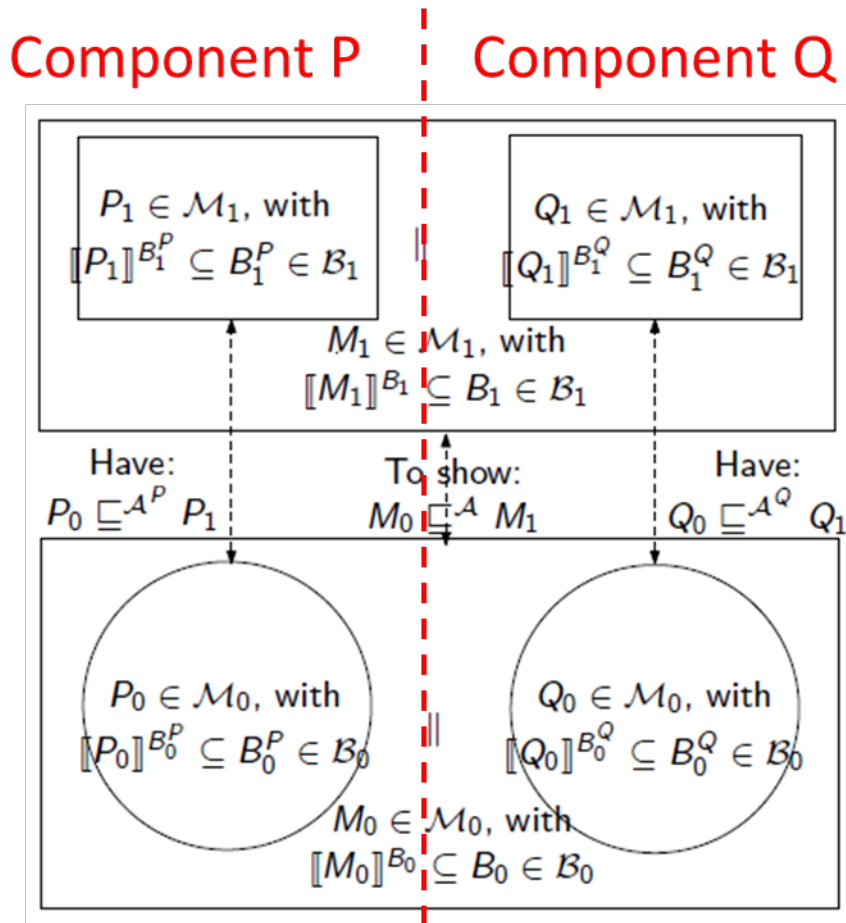
If $M_0 \sqsubseteq^{R_1} M_1$, $M_1 \models^{B_1} S_1$ and $S_1 \Rightarrow^{R_1} S_0$,
then $M_0 \models^{B_0} S_0$.

heterogeneous
abstraction

M_1	\models^{B_1}	S_1
\sqsubseteq^{R_1}		\Rightarrow^{R_1}
M_0	\models^{B_0}	S_0

heterogeneous
implication

Compositional Heterogeneous Abstraction



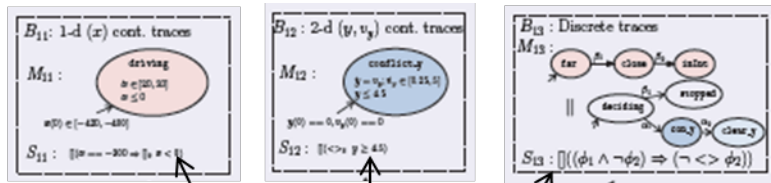
When is the composition of component abstractions an abstraction of the composition of components?

Sufficient Condition

Abstraction relations $\mathcal{A}^P, \mathcal{A}^Q$ have a common globalization \mathcal{A} .

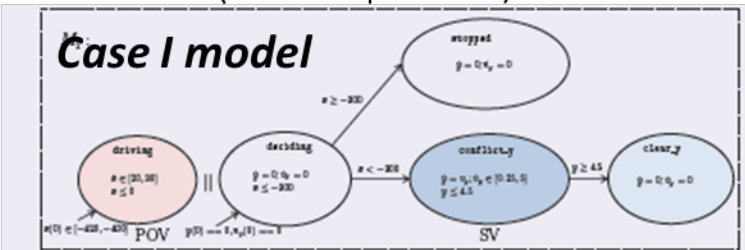
Note: A single class of behaviors is used at each level.

Application to CICAS-SSA



Conjunctive breakdowns

- M_{1i} individually cover M_1
- S_{1i} together imply S_1

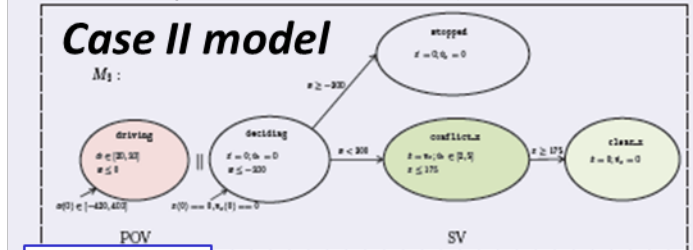


No right turn

B_1 : 3-d ($\{x, y, v_y\}$) hybrid traces
 S_1 : $\square \neg (x == 0 \wedge 0 < y < 4.5)$

R_1 : equal values along $x, y, v_y; z, v_z$ zero over all time

$R_1^{-1}([S_1]^{B_1})$: $\square ((\neg(x == 0 \wedge 0 < y < 4.5)) \wedge (z == 0 \wedge v_z == 0))$



No straight crossing

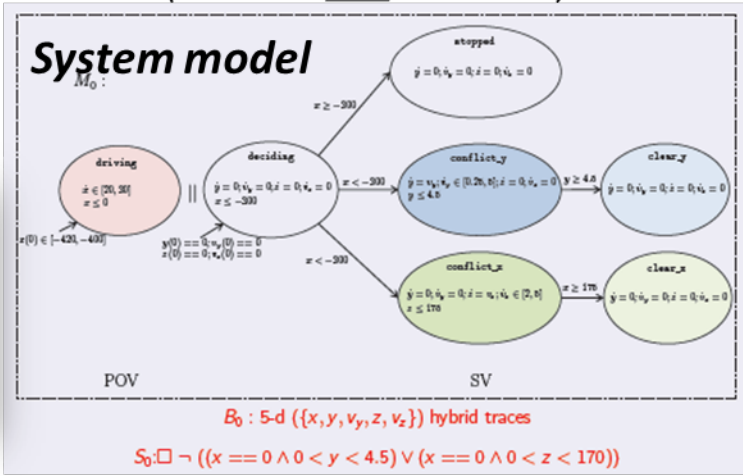
B_2 : 3-d ($\{x, z, v_z\}$) hybrid traces
 S_2 : $\square \neg (x == 0 \wedge 0 < z < 170)$

R_2 : equal values along $x, z, v_z; y, v_y$ zero over all time

$R_2^{-1}([S_2]^{B_2})$: $\square ((\neg(x == 0 \wedge 0 < z < 170)) \wedge (y == 0 \wedge v_y == 0))$

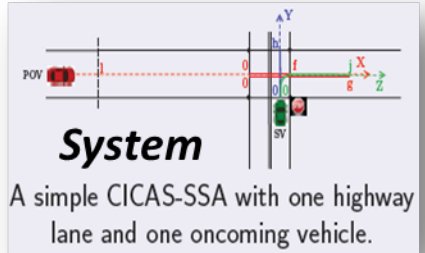
Disjunctive breakdown

- M_i together cover M_0
- S_i individually imply S_0



B_0 : 5-d ($\{x, y, v_y, z, v_z\}$) hybrid traces
 S_0 : $\square \neg ((x == 0 \wedge 0 < y < 4.5) \vee (x == 0 \wedge 0 < z < 170))$

★ 5-d hybrid verification simplified to 1-d, 2-d continuous and discrete verifications



Heterogeneous Verification via Behavioral Semantics

- Behavior domain **relations** support heterogeneous
 - abstraction
 - implication
 - compositional abstraction
- Applications to CICAS-SSA
- Formalizes specific cases in the literature
 - papers: **google** “Rajhans Krogh heterogeneous”

Using Architecture to Manage Formal Analysis

